

Parallel Temporal Decomposition for Improved Unit Commitment in Power System Production Cost Modeling

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GMLC: Multi-Scale Production Cost Models

- GMLC: Grid Modernization Laboratory Consortium
 - An aggressive five-year grid modernization strategy for the Department of Energy
- Design and planning tools sub-area includes Multi-Scale Production Cost Models
 - Develop multi-scale production cost models with faster mathematical solvers
- PCM Goal:
 - Substantially increase the ability of production cost models (PCM) to simulate power systems in more detail faster and more robustly.
 - Both Deterministic and Stochastic
- Talks at Technical Conference:
 - Session T1-B: Optimization Driven Scenario Grouping for Stochastic Unit Commitment (LLNL)
 - Session T2-B: Assessment of Wind Power Ramp Events in Scenario Generation for Stochastic Unit Commitment (SNL)
 - Session T3-A: Geographic Decomposition of Production Cost Models (NREL)
 - Session T3-A: Temporal Decomposition of the Production Cost Modeling in Power Systems (ANL)









Power System Operations





PCM: Unit Commitment and Economic Dispatch



- Unit Commitment: scheduling generators on/off
- Economic Dispatch: scheduling power generation at each generator
- Security Constraints:
 - Flow balance constraints
 - Power flow constraints
 - Ramping constraints
 - Minimum up/down constraints
 - Spinning reserve constraints





WECC System

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Multi-Scale Production Cost Modeling (PCM)

Goal

to simulate a broad range of scenarios in order to plan electricity system over a long-term planning horizon

Challenges

- The complexity and resolution required to model the modern power system is rapidly increasing.
- Model fidelity vs. execution time
- Needs to solve long-term unit commitment and economic dispatch







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Temporal Decomposition





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Long-Term UC Model Formulation

$$\begin{array}{ll} \min & \sum_{g \in G} \sum_{t \in T} \left(K_g u_{gt} + S_g v_{gt} + C_g p_{gt} \right) & \text{Operating cost} \\ \text{s.t.} & \sum_{l \in L_n^+} f_{lt} - \sum_{l \in L_n^-} f_{lt} + \sum_{g \in G_n} p_{gt} = D_{nt}, \ n \in N, \ t \in T, & \text{Flow balance equation} \\ f_{ll} = B_l \left(\theta_{nt} - \theta_{mt} \right), \ l = (m, n) \in L, \ t \in T, & \text{Linearized power flow equation} \\ & - F_l \leq f_{lt} \leq F_l, \ l \in L, \ t \in T, & \text{Transmission line capacity} \\ s_{gt} \leq p_{gt} \leq r_{gt}, \ g \in G, \ t \in T, & \text{Operating reserve requirement} \\ r_{gt} \leq P_g^{max} u_{gt}, \ g \in G, \ t \in T, & \text{Generation capacity} \\ s_{gt} \geq P_g^{min} u_{gt}, \ g \in G, \ t \geq T, & \text{Generation capacity} \\ s_{gt} - p_{g,t-1} \leq R_g^+, \ g \in G, \ t \geq 2 & \text{Ramping capacity} \\ s_{gt} - p_{g,t-1} \geq -R_g^-, \ g \in G, \ t \geq 2 & \text{Ninimum uptime} \\ downtime requirements \\ \end{array}$$



Temporal Decomposition – Linking Constraints

$$\begin{array}{ll} \min & \sum_{g \in G} \sum_{t \in T} \left(K_g u_{gt} + S_g v_{gt} + C_g p_{gt} \right) \\ \text{s.t.} & \sum_{l \in L_n^+} f_{lt} - \sum_{l \in L_n^-} f_{lt} + \sum_{g \in G_n} p_{gt} = D_{nt}, \ n \in N, \ t \in I \\ & f_{lt} = B_l \left(\theta_{nt} - \theta_{mt} \right), \ l = (m, n) \in L, \ t \in T, \\ & - F_l \leq f_{lt} \leq F_l, \ l \in L, \ t \in T, \\ & s_{gt} \leq p_{gt} \leq r_{gt}, \ g \in G, \ t \in T, \\ & r_{gt} \leq P_g^{max} u_{gt}, \ g \in G, \ t \in T, \\ & s_{gt} \geq P_g^{min} u_{gt}, \ g \in G, \ t \in T, \\ & r_{gt} - p_{g,t-1} \leq R_g^+, \ g \in G, \ t \geq 2 \\ & s_{gt} - p_{g,t-1} \geq -R_g^-, \ g \in G, \ t \geq 2 \\ & \sum_{q=t-UT_g+1}^t v_{gq} \leq u_{gt}, \ g \in G, \ t \geq UT_g, \\ & v_{gt} - w_{gt} = u_{gt} - u_{g,t-1}, \ g \in G, \ t \geq 2, \\ & u_{gt}, v_{gt}, w_{gt} \in \{0,1\}, \ g \in G, \ t \in T \end{array}$$



T,

Coupling Constraints that link the variables in different time periods



Temporal Decomposition – Linking Constraints

Subintervals of the time horizon:

- 1. $T_j \subset T$ for $j \in J$;
- 2. $\cup_{j \in J} T_j = T$; and
- 3. $T_i \cap T_j = \emptyset$ for $i \neq j \in J$.

Equivalent PCM formulation:

$$\min \sum_{g \in G} \sum_{t \in T} (K_g u_{gt} + S_g v_{gt} + C_g p_{gt})$$
s.t. $r_{gt} - p_{g,t-1} \leq R_g^+, \ g \in G, \ t \in T_j, \ t - 1 \notin T_j, \ j \in J,$
 $s_{gt} - p_{g,t-1} \geq -R_g^-, \ g \in G, \ t \in T_j, \ t - 1 \notin T_j, \ j \in J,$
 $\sum_{q=t-UT_g+1}^t v_{gq} \leq u_{gt}, \ g \in G, \ t \in T_j, \ t - UT_g + 1 \notin T_j, \ j \in J,$
 $v_{gt} - w_{gt} = u_{gt} - u_{g,t-1}, \ g \in G, \ t \in T_j, \ t - DT_g + 1 \notin T_j, \ j \in J,$
 $v_{gt} - w_{gt} = u_{gt} - u_{g,t-1}, \ g \in G, \ t \in T_j, \ t - 1 \notin T_j, \ j \in J,$
 $(\mathbf{u}_j, \mathbf{v}_j, \mathbf{w}_j, \mathbf{p}_j, \mathbf{r}_j, \mathbf{s}_j) \in \mathcal{X}_j, \ j \in J,$
 $mixed-binary set$



We simplify the formulation.

min
$$\sum_{j \in J} \sum_{k \in K} \mathbf{c}_j \mathbf{x}_j$$

s.t.
$$\sum_{j \in J} \sum_{k \in K} \mathbf{A}_j \mathbf{x}_j \ge \mathbf{b},$$

$$\mathbf{x}_j \in conv(\mathcal{X}_j), \ j \in J$$



Decomposition Methods

Simplified Original Formulation:

$$\min \quad \sum_{j \in J} \sum_{k \in K} \mathbf{c}_j \mathbf{x}_j \\ \text{s.t.} \quad \sum_{j \in J} \sum_{k \in K} \mathbf{A}_j \mathbf{x}_j \ge \mathbf{b}, \\ \mathbf{x}_j \in conv(\mathcal{X}_j), \ j \in J$$

Dantzig-Wolfe Decomposition

min
$$\sum_{j \in J} \sum_{k \in K} \mathbf{c}_j \mathbf{x}_j^k \alpha_j^k$$

s.t.
$$\sum_{j \in J} \sum_{k \in K} \mathbf{A}_j \mathbf{x}_j^k \alpha_j^k \ge \mathbf{b} \ (\lambda),$$

$$\sum_{k \in K} \alpha_j^k = 1 \ (\mu_j), \ j \in J$$
$$\alpha_j^k \ge 0, \ j \in J, \ k \in K,$$

Dual Decomposition

$$\max_{\lambda \ge 0} \min \sum_{j \in J} \mathbf{c}_j \mathbf{x}_j + \lambda^T \left(\mathbf{b} - \sum_{j \in J} \mathbf{A}_j \mathbf{x}_j \right)$$

s.t. $\mathbf{x}_j \in conv(\mathcal{X}_j), \ j \in J$

Dual Decomposition

- Outer approximation
- Row generation procedure
- Dual solution space
- Regularization

Dantzig-Wolfe Decomposition

- Inner approximation
- Column generation procedure
- Primal solution space
- Branching, heuristics



Dual Decomposition

Simplified Original Formulation:

$$\min \sum_{j \in J} \sum_{k \in K} \mathbf{c}_j \mathbf{x}_j$$
s.t.
$$\sum_{j \in J} \sum_{k \in K} \mathbf{A}_j \mathbf{x}_j \ge \mathbf{b},$$

$$\mathbf{x}_j \in conv(\mathcal{X}_j), \ j \in J$$

Dual Decomposition

$$\max_{\lambda \ge 0} \min \sum_{j \in J} \mathbf{c}_j \mathbf{x}_j + \lambda^T \left(\mathbf{b} - \sum_{j \in J} \mathbf{A}_j \mathbf{x}_j \right)$$

s.t. $\mathbf{x}_j \in conv(\mathcal{X}_j), \ j \in J$



Maximizing the Lagrangian dual bound

$\max_{\lambda \ge 0} \left\{ \sum_{j \in J} D_j(\lambda) + \mathbf{b}^T \lambda \right\} \qquad \begin{array}{l} \text{Decomposed for each } j.\\ D_j(\lambda) := \min\left\{ \left(\mathbf{c}_j - \lambda^T \mathbf{A}_j \right) \mathbf{x}_j : \mathbf{x}_j \in conv(\mathcal{X}_j) \right\} \end{array}$

Proximal bundle model:

$$\begin{aligned} \max \quad & \sum_{j \in J} \mu_j + \mathbf{b}^T \lambda - \frac{1}{2} \|\lambda - \hat{\lambda}\|_2^2 \\ \text{s.t.} \quad & \mu_j \leq D_j (\lambda^k) - (\mathbf{A}_j \mathbf{x}_j)^T (\lambda - \lambda^k), \ j \in J, \ k \in K, \\ & \lambda \geq 0 \end{aligned}$$



Dantzig-Wolfe Decomposition

Simplified Original Formulation:

 $\begin{array}{ll} \min & \sum_{j \in J} \sum_{k \in K} \mathbf{c}_j \mathbf{x}_j \\ \text{s.t.} & \sum_{j \in J} \sum_{k \in K} \mathbf{A}_j \mathbf{x}_j \geq \mathbf{b}, \\ & \mathbf{x}_j \in conv(\mathcal{X}_j), \ j \in J \end{array}$

Dantzig-Wolfe Decomposition

$$\min \sum_{j \in J} \sum_{k \in K} \mathbf{c}_j \mathbf{x}_j^k \alpha_j^k$$

s.t.
$$\sum_{j \in J} \sum_{k \in K} \mathbf{A}_j \mathbf{x}_j^k \alpha_j^k \ge \mathbf{b} \ (\lambda),$$
$$\sum_{k \in K} \alpha_j^k = 1 \ (\mu_j), \ j \in J$$
$$\alpha_j^k \ge 0, \ j \in J, \ k \in K,$$

Pricing Problems: Add new column if

$$\mu_j^k > \min\left\{ \left(\mathbf{c}_j - (\lambda^k)^T \mathbf{A}_j \right) \mathbf{x}_j : \mathbf{x}_j \in conv(\mathcal{X}_j) \right\}$$

- Decomposed for each j
- Same as the *dual decomposition subproblems*

Valid Lagrangian Dual Bound:

$$z_{DW} \ge \sum_{j \in J} \left(\mathbf{c}_j - (\lambda^k)^T \mathbf{A}_j \right) \mathbf{x}_j^k + \mathbf{b}^T \lambda^k$$

Best Lagrangian Dual Bound:

$$z \geq z_{DW} = \sum_{j \in J} \left(\mathbf{c}_j - (\lambda^k)^T \mathbf{A}_j \right) \mathbf{x}_j^k + \mathbf{b}^T \lambda^k$$

Positive duality gap may exist.



Integrating Branch-and-Bound Method

Simplified Original Formulation:

$$\min \quad \sum_{j \in J} \sum_{k \in K} \mathbf{c}_{j} \mathbf{x}_{j} \\ \text{s.t.} \quad \sum_{j \in J} \sum_{k \in K} \mathbf{A}_{j} \mathbf{x}_{j} \ge \mathbf{b}, \\ \mathbf{x}_{j} \in conv(\mathcal{X}_{j}), \ j \in J$$

Recovering Original Variables:

$$\mathbf{x}_j = \sum_{k \in K} \mathbf{x}_j^k \alpha_j^k \in conv(\mathcal{X}_j)$$

Convex combination of feasible solutions

Not necessarily integer feasible

Branching in DW Decomposition:



Adding branching rows

Also adding branching columns to the DD master problem



Dantzig-Wolfe Decomposition

 $\min \sum_{j \in J} \sum_{k \in K} \mathbf{c}_j \mathbf{x}_j^k \alpha_j^k$ s.t. $\sum_{j \in J} \sum_{k \in K} \mathbf{A}_j \mathbf{x}_j^k \alpha_j^k \ge \mathbf{b} \ (\lambda),$ $\sum_{k \in K} \alpha_j^k = 1 \ (\mu_j), \ j \in J$ $\alpha_j^k \ge 0, \ j \in J, \ k \in K,$

Flowchart – Temporal Decomposition





DSP: Scalable Decomposition Solver

Decomposition methods for Structured Programming

- Exploiting block-angular structures
- Dantzig-Wolfe decomposition
 + (Parallel) Branch-and-Bound
- Benders decomposition
- Dual decompositioParallel
- Parallel computing via MPI library

| C This re | epository Search | | Pull requests | Issues Gi | st | | | 🖍 +- 🐉 | |
|-----------------------------------|------------------|--------------------|----------------|-----------|-----------|----------------|------------|-------------|--|
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| <> Code | () Issues () | 11 Pull requests 0 | III Projects 0 | 🗉 Wiki | -/~ Pulse | III Graphs | | | |

An open-source parallel optimization solver for stochastic mixed-integer programming

| 299 commits | ₽ 5 branches | ℃ 5 releases | | 4 contributors | |
|-----------------------------------|--|-----------------|--------------|----------------|-------------------------|
| Branch: master - New pull request | | Create new file | Upload files | Find file | Clone or download - |
| Kibaek Kim bugfix #11 | | | | Latest comr | nit edc6c0b 18 days ago |
| docs | new version release | | | | 9 months ago |
| examples | bugfix #11 | | | | 18 days ago |
| extra | fix ambiguous definition of hash for gcc 6 | | | | 8 months ago |
| iii lib | untrack libDsp.so | | | | a year ago |
| parameters | new version release | | | | 9 months ago |
| src src | bugfix #11 | | | | 18 days ago |









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DSP reads models from Julia



where

$$Q(x_1, x_2, \xi_1^s, \xi_2^s) = \min_{y_1, y_2, y_3, y_4} -16y_1 + 19y_2 + 23y_3 + 28y_4$$

s.t.
$$2y_1 + 3y_2 + 4y_3 + 5y_4 \le \xi_1^s - x_1$$
$$6y_1 + y_2 + 3y_3 + 2y_4 \le \xi_2^s - x_2$$
$$y_1, y_2, y_3, y_4 \in \{0, 1\}$$

and $(\xi_1^s, \xi_2^s) \in \{(7, 7), (11, 11), (13, 13)\}$ with probability 1/3.

```
1 using Dsp, MPI # Load packages
```

```
2 MPI.Init() # Initialize MPI
```

Only 15 lines of Julia script!

```
3 m = Model(3) # Create a Model object with three scenarios
```

```
4 xi = [[7,7] [11,11] [13,13]] # random parameter
```

```
5 @variable(m, 0 <= x[i=1:2] <= 5, Int)
```

```
6 @objective(m, Min, -1.5*x[1]-4*x[2])
```

```
7 for s in 1:3
```

julia

```
8 q = Model(m, s, 1/3);
```

```
9 @variable(q, y[j=1:4], Bin)
```

```
10 @objective(q, Min, -16*y[1]+19*y[2]+23*y[3]+28*y[4])
```

```
11 @constraint(q, 2*y[1]+3*y[2]+4*y[3]+5*y[4]<=xi[1,s]-x[1])
```

```
12 @constraint(q, 6*y[1]+1*y[2]+3*y[3]+2*y[4]<=xi[2,s]-x[2])
```

13 **end**

```
14 solve(m, solve_type=:Dual, param="myparams.txt")
```

```
15 MPI.Finalize() # Finalize MPI
```



Computational Results – IEEE 118-Bus System

Implementation

- Argonne's parallel open-source software: DSP + Coin-ALPS
- Julia interface for modeling
- Running on Argonne's Blues cluster (600-node computing node with 16 cores on each)

IEEE 118-Bus System

 118 buses, 54 generators, and 186 transmission lines
 Estimated hourly demand from the PJM system (April 2016)

Table 1: Sizes of IEEE 118-bus system problem instances

| | 1 | # Constraints | # Variables | # Binary |
|---|-----|---------------|-------------|----------|
| - | 24 | 19765 | 18960 | 1296 |
| | 48 | 40070 | 37920 | 2592 |
| | 72 | 60398 | 56880 | 3888 |
| | 96 | 80726 | 75840 | 5184 |
| | 120 | 101054 | 94800 | 6480 |
| | 144 | 121382 | 113760 | 7776 |
| | 168 | 141710 | 132720 | 9072 |





Solution Time Benchmark – Log Scale





Computational Results - Branch-and-Bound





Computational Results - Branch-and-Bound





Computational Results - Branch-and-Bound





Future Work

Algorithms

- Integrating with generic MIP solution
- Inexact subproblem solutions
- Primal cutting planes
- Primal heuristics
- Further Parallelization
 - Master problems
 - Tree search
- Investigating Applications
 - Network decomposition
 - Hybrid decomposition (network, time, scenarios)
 - Other applications

